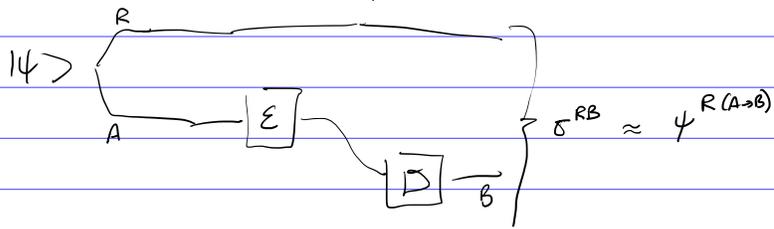


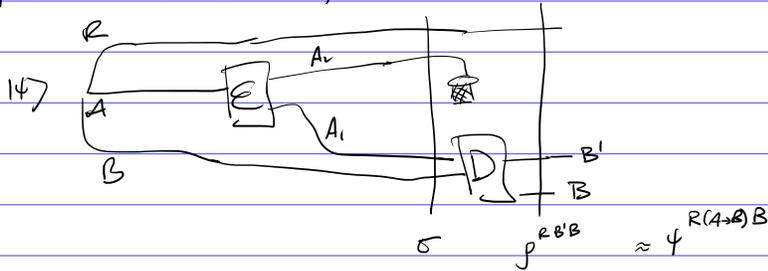
Merging: application of random unitaries to q. info. theory

recd Schumacher compression



A to B means it's controlled by B
but looks like the original A system

compression with side information



Strategy

$$A = A_1 \otimes A_2 \quad d_A = d_{A_1} d_{A_2}$$

E = random unitary $U^{A \rightarrow A_1 A_2} \sim \text{Haar}$ (or 2-design)

decoupling: $\sigma^{A_2 R} \approx \sigma^{A_2} \otimes \sigma^R$ (*)

purify \downarrow
 $(U \otimes I_{BR}) |\psi\rangle \rightarrow |\psi\rangle^{A_2 \tilde{B}} \otimes |\psi\rangle^{R B' B}$

Uhlmann $\Rightarrow \exists V: A, B \rightarrow \tilde{B} B' B$ s.t.

$$V_{A, B} |\psi\rangle \approx |\psi\rangle^{A_2 \tilde{B}} |\psi\rangle^{R B' B}$$



Pf of (*)

First calculate $\mathbb{E} \|\sigma^{A_2 R} - \sigma^{A_2} \otimes \sigma^R\|_2^2 = \mathbb{E} [\text{tr} \sigma_{A_2 R}^2 - 2 \text{tr} \sigma_{A_2 R} (\sigma_{A_2} \otimes \sigma^R) + \text{tr} \sigma_{A_2}^2 \text{tr} \sigma^R]$

$$\begin{aligned} \mathbb{E} \text{tr} \sigma_{A_2 R}^2 &= \mathbb{E} \text{tr} (\sigma_{A_2 R} \otimes \sigma_{A_2 R}^T) F_{A_2} F_R \\ &= \mathbb{E} \text{tr} (U \psi_{AR} U^\dagger \otimes U \psi_{AR} U^\dagger) F_{A_2} F_R \\ &= \text{tr} (\psi_{AR} \otimes \psi_{AR}^T) \mathbb{E} (U^\dagger \otimes U^\dagger) F_{A_2} (U \otimes U) F_R \\ &= \alpha_+ \Pi_+ + \alpha_- \Pi_- \end{aligned}$$

$$\begin{aligned} \Pi_\pm &= \frac{I \pm F_A F_{A_2}}{2} & \alpha_\pm &= \frac{\text{tr} \Pi_\pm F_{A_2}}{\text{tr} F_{A_2}} \\ & & &= \frac{1}{d_{A_1} \pm d_{A_1}} \\ & & &= d_{A_1}^{-1} d_{A_2} \pm d_{A_1} d_{A_2}^{-1} = d_A (d_{A_1} \pm d_{A_2}) \end{aligned}$$

$$\begin{aligned} &= \frac{\alpha_+ + \alpha_-}{2} I + \frac{\alpha_+ - \alpha_-}{2} F_A \approx d_{A_2}^{-1} I + d_{A_1}^{-1} F_A \\ \Rightarrow \mathbb{E} \text{tr} \sigma_{A_2 R}^2 &\approx d_{A_2}^{-1} \text{tr} \psi_R^2 + d_{A_1}^{-1} \text{tr} \psi_{AR} \end{aligned}$$

$$\mathbb{E} \operatorname{tr} \sigma_{A_2}^2 + \operatorname{tr} \sigma_R^2 = \mathbb{E} (\operatorname{tr} \sigma_{A_2}^2) + \operatorname{tr} \psi_R^2$$

$$\mathbb{E} \operatorname{tr} \sigma_{A_2}^2 = \operatorname{tr} \psi_{AR} \otimes \psi_{AR} \mathbb{E} (U^\dagger \otimes U) F_{A_2} (U \otimes U) \approx d_{A_2}^{-1} + d_{A_1}^{-1} \operatorname{tr} \psi_A^2 \approx d_{A_2}^{-1}$$

if $d_{A_1} > d_{A_2}$

$$\begin{aligned} \mathbb{E} \operatorname{tr} \sigma_{AR} (\sigma_{A_2} \otimes \sigma_R) &= \mathbb{E} \operatorname{tr} (U \otimes U) (\psi_{AR} \otimes \psi_A \otimes \psi_R) (U^\dagger \otimes U) F_A F_R \\ &= \operatorname{tr} (\psi_{AR} \otimes \psi_A \otimes \psi_R) (d_{A_2}^{-1} F_R + d_{A_1}^{-1} F_A F_R) \\ &= d_{A_2}^{-1} \operatorname{tr} \psi_R^2 + d_{A_1}^{-1} \operatorname{tr} \psi_{AR} (\psi_A \otimes \psi_R) \end{aligned}$$

$d_{A_2}^{-1} \operatorname{tr} \psi_R^2$ terms cancel

$$\Rightarrow \mathbb{E} \|\sigma_{AR} - \sigma_{A_2} \otimes \sigma_R\|_2^2 \approx d_{A_1}^{-1} (\operatorname{tr} \psi_{AR}^2 - 2 \operatorname{tr} \psi_{AR} (\psi_A \otimes \psi_R) + \operatorname{tr} \psi_A^2 \operatorname{tr} \psi_R^2)$$

$\exists U$ s.t. $\|\cdot\|_2^2 \leq \dots$

Fix R is U .

$$\|\sigma_{AR} - \sigma_{A_2} \otimes \sigma_R\|_1^2 \leq \frac{d_{A_2} d_R}{d_{A_1}} (\operatorname{tr} \psi_{AR}^2 + \operatorname{tr} \psi_A^2 \operatorname{tr} \psi_R^2)$$

Is this small? yes for $\frac{d_{A_1}}{d_{A_2}}$ big enough.

concretely

$$\rho_{AB} \xrightarrow{\text{purify}} |\phi\rangle_{ABR} \rightarrow |\psi\rangle = c (\prod_{\phi_{AS}}^n \otimes \prod_{\psi_{BS}}^n \otimes \prod_{\psi_{RS}}^n) |\phi\rangle^{\otimes n}$$

$$\operatorname{tr} \psi_A^2 \approx \exp(-n S(A)_\rho) = \exp(-n S(A)_p)$$

$$\operatorname{tr} \psi_R^2 \approx \exp(-n S(R)_\rho) = \exp(-n S(AB)_p)$$

$$\operatorname{tr} \psi_{AR}^2 \approx \exp(-n S(AR)_\rho) = \exp(-n S(B)_p) \approx \exp(-n (S(A) + S(R)))$$

$$d_R \approx \exp(n S(AB)_p) \quad d_A \approx \exp(n S(A)_p)$$

error small if $d_A d_R \operatorname{tr} \psi_{AR}^2 \ll d_{A_1}^2$

$$\log d_{A_1} \gg n \frac{I(A;R)_p}{2}$$

$$\text{yields } \log d_{A_2} \text{ ebits} \quad \log d_{A_2} \approx n \left(S(A) - \frac{I(A;R)}{2} \right) = n \frac{I(A;B)}{2}$$

$$\langle \phi^{ABR} \rangle + \frac{I(A;R)}{2} \text{ qubits} \geq \langle \phi^{(A \rightarrow B)BR} \rangle + \frac{I(A;B)}{2} \text{ ebits}$$

If cbits are free then the entanglement cost is

$$\begin{aligned} \frac{I(A;R) - I(A;B)}{2} &= \frac{1}{2} (S(A) + S(R) - S(R) - (S(A) + S(B) - S(R))) = S(R) - S(B) \\ &= S(AB) - S(B) = S(A|B) \end{aligned}$$